

# Pseudofinite countably categorical theories

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## Definition

A theory  $T$  is *pseudofinite* if every sentence  $\varphi \in T$  has a finite model. Equivalently, some model of  $T$  is an ultraproduct of finite structures.

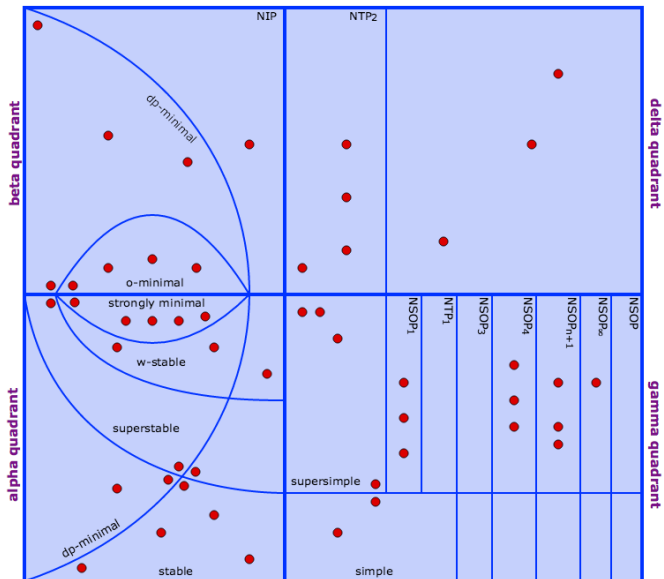
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## Question

Where do  $\aleph_0$ -categorical pseudofinite theories lie relative to model-theoretic dividing lines?

# Map of the universe (Gabe Conant)



## Theorem (Folklore)

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By compactness, we can find a densely ordered chain  $\{a_i \mid i \in [0, 1] \cap \mathbb{Q}\}$  with  $a_i < a_j \leftrightarrow i < j$ . By  $\aleph_0$ -categoricity, there is some formula  $\varphi(x, y)$  such that  $P \models \varphi(b, c)$  if and only if there is a densely ordered chain between  $b$  and  $c$ .

# Pseudofiniteness and SOP

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Now  $P \models (\exists x \exists y \varphi(x, y)) \wedge (\forall x \forall y \varphi(x, y) \rightarrow \exists z (x < z \wedge \varphi(z, y)))$ .

This sentence cannot hold in any finite structure - it implies the existence of an infinite ascending chain. □



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Quasi-finitely axiomatizable: Theory of vector spaces +  $\{\exists^{\geq n} x \top(x) \mid n \in \omega\}$ .

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This example is one of a large family, the smoothly approximable structures, classified by Cherlin and Hrushovski (*Finite Structures with Few Types*). Built from pure sets and geometries coming from vector spaces, possibly with bilinear forms.

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*Example:* The random graph.

Axiomatizable by extension axioms: Theory of graphs +  
 $\{\forall \bar{x} \forall \bar{y} \exists z (\bigwedge_{i,j} x_i \neq y_j) \rightarrow (\bigwedge_{i=1}^n z E x_i) \wedge (\bigwedge_{j=1}^n \neg z E y_j) \mid n \in \omega\}$ .

Let  $\mathcal{G}(n)$  be the graphs with domain  $[n] = \{1, \dots, n\}$ , and take  $\mu_n$  to be the uniform measure on  $\mathcal{G}(n)$ . Then for any sentence  $\varphi$  in the theory of the random graph,  $\lim_{n \rightarrow \infty} \mu_n(\{G \in \mathcal{G}(n) \mid G \models \varphi\}) = 1$ .

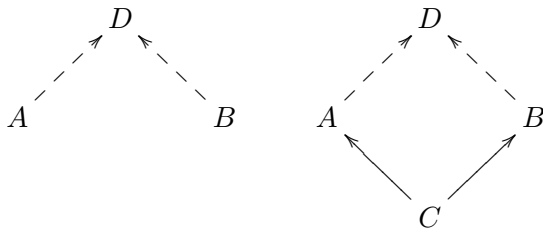
We say that the theory of the random graph is the *almost-sure theory* of the sequence  $(\mathcal{G}(n), \mu_n)_{n \in \omega}$ .

So every sentence of this theory has a rich class of finite models: Almost all large finite graphs.

# A brief review of Fraïssé theory

Let  $L$  be a relational language. A *Fraïssé class* is a class  $K$  of finite  $L$ -structures, closed under isomorphism, such that

- 1  $K$  is countable up to isomorphism.
- 2  $K$  has the hereditary property: If  $B \in K$  and  $A$  embeds in  $B$ , then  $A \in K$ .
- 3  $K$  has the joint embedding property and the amalgamation property:



Since we are interested in countably categorical theories, we assume that our Fraïssé classes contain only finitely many structures of size  $n$  up to isomorphism for each  $n$ .

# A brief review of Fraïssé theory

Let  $K$  be a Fraïssé class. Then there is a unique countable structure up to isomorphism  $M_K$ , the *Fraïssé limit* of  $K$ , such that

- 1  $K$  is the class of finite substructures of  $M_K$ .
- 2  $M_K$  is homogeneous: Any isomorphism between finite substructures extends to an automorphism of  $M_K$ .

Let  $T_K = \text{Th}(M_K)$ . Then  $T_K$  is  $\aleph_0$ -categorical, has quantifier elimination, and can be explicitly axiomatized by:

- 1 Universal axioms asserting that every finite substructure is in  $K$ .
- 2 A  $\forall\exists$  *extension axiom* for every embedding  $A \rightarrow B$  in  $K$ :

$$\forall \bar{x} \exists \bar{y} \theta_A(\bar{x}) \rightarrow \theta_B(\bar{x}, \bar{y})$$

where  $\theta_X$  describes the quantifier-free type of  $X$ .

$T_K$  is called the *generic theory* of  $K$ .

# The canonical language

Let  $T$  be any  $\aleph_0$ -categorical theory and  $M$  its unique countable model.

- Let  $L^*$  be the language containing one  $n$ -ary predicate for each  $n$ -type relative to  $T$ .  $L^*$  is called the *canonical language* for  $T$ .
- Give  $M^*$  is natural  $L^*$ -structure.
- The class  $K_{T^*}$  of all finite substructures of  $M^*$  is a Fraïssé class, with Fraïssé limit  $M^*$  and generic theory  $T^*$ .

$T$  and  $T^*$  are *interdefinable*.



## Definition

Let  $K$  and  $K'$  be Fraïssé classes in languages  $L$  and  $L'$ , respectively, such that  $L \subseteq L'$ . We say that  $K$  is an expansion of  $K'$  if

- 1  $K = \{A \upharpoonright L \mid A \in K'\}$
- 2 For all extensions  $(A, B)$  in  $K$ , and every expansion of  $A$  to a structure  $A'$  in  $K'$ , there is an expansion of  $B$  to a structure  $B'$  in  $K'$  such that  $(A', B')$  is an extension.

$$\begin{array}{ccc} A & \longrightarrow & B \\ \downarrow \upharpoonright L & & \downarrow \upharpoonright L \\ A' & \dashrightarrow & B' \end{array}$$

$K'$  is an expansion of  $K$  if and only if the Fraïssé limit  $M_{K'}$  of  $K'$  is an expansion of the Fraïssé limit  $M_K$  of  $K$ .

# Disjoint $n$ -amalgamation

## Theorem (K.)

*If a countably categorical theory  $T$  has disjoint  $n$ -amalgamation for all  $n$ , then  $T$  is pseudofinite. Consequently, any reduct of  $T$  is also pseudofinite.*

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A type  $p(\bar{x})$  over  $A$  is *non-redundant* if it does not contain the formulas  $x_i = x_j$  for  $i \neq j$  or  $x_i = a$  for  $a \in A$ .

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A theory has *disjoint  $n$ -amalgamation* if whenever we have tuples of variables  $\bar{x}_1, \dots, \bar{x}_n$  and a system of types  $\{p_S \mid S \subsetneq [n]\}$  over  $A$  such that  $p_S$  is a non-redundant type in the variables  $\{\bar{x}_i \mid i \in S\}$ , and  $p_S \subseteq p_T$  when  $S \subseteq T$ , then there is some non-redundant type  $p_{[n]}$  in the variables  $\{\bar{x}_i \mid i \in [n]\}$  extending the  $p_S$ .

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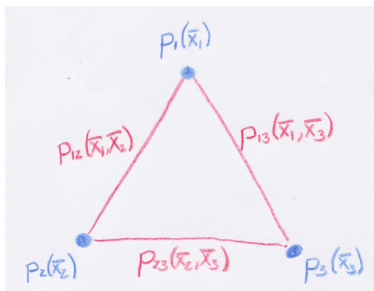
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## Fact

*A countably categorical theory has disjoint 2-amalgamation if and only if it has trivial algebraic closure:  $\text{acl}(A) = A$  for all sets  $A$ .*

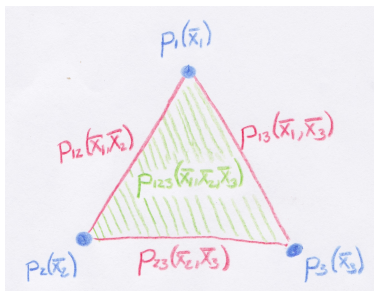
# 3-amalgamation

Given  $p_i(\bar{x}_i)$  and  $p_{ij}(\bar{x}_i, \bar{x}_j)$ ,  $i, j \in \{1, 2, 3\}$ , can we find  $p_{123}(\bar{x}_1, \bar{x}_2, \bar{x}_3)$ ?



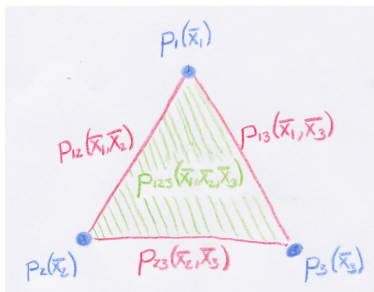
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**Examples of failure:** Let  $T_K$  be the generic theory of a Fraïssé class  $K$ .

- $K =$  partial orders.  $x_1 < x_2$ ,  $x_2 < x_3$ ,  $x_3 < x_1$ .
- $K =$  equivalence relations.  $x_1 E x_2$ ,  $x_2 E x_3$ ,  $\neg x_1 E x_3$ .
- $K =$  triangle-free graphs.  $x_1 E x_2$ ,  $x_2 E x_3$ ,  $x_1 E x_3$ .



## Theorem (K.)

*If a countably categorical theory  $T$  has disjoint  $n$ -amalgamation for all  $n$ , then  $T$  is pseudofinite. Consequently, any reduct of  $T$  is also pseudofinite.*

Using the canonical language, we may assume that  $T$  is the generic theory of a Fraïssé class  $K$ . We describe a probability measure  $\mu_n$  on  $K(n)$ , the structures in  $K$  with domain  $[n] = \{1, \dots, n\}$  by giving a probabilistic construction of such a structure. Then we check that  $T$  is the almost-sure theory of  $(K(n), \mu_n)_{n \in \omega}$ .

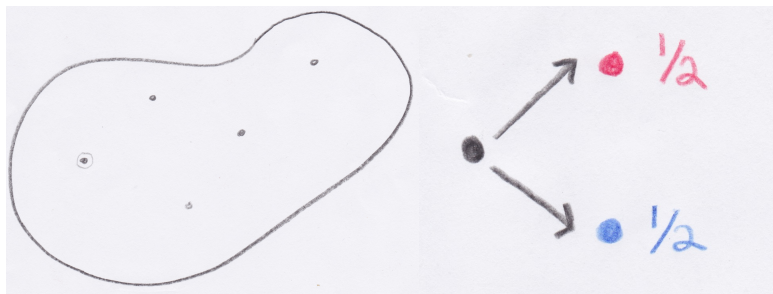
The measure  $\mu_n$  is typically *not* the uniform measure on  $K(n)$ , but it is still very *natural*, and the probabilistic argument produces a large number of finite models in  $K$  for every sentence in  $T$ .

# Probabilistic construction

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We build from below, assigning quantifier-free 1-types to each element, uniformly at random.



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Then assign quantifier-free 2-types to each pair, uniformly among those extending the given 1-types.

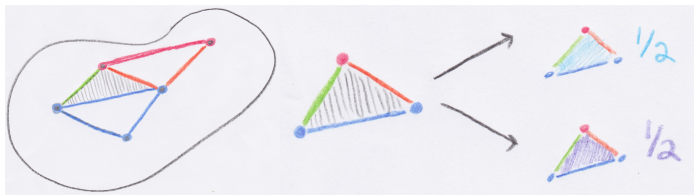


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Disjoint  $n$ -amalgamation for all  $n$  ensures that we never get stuck.



... and that all choices were made as independently as possible. This allows us to verify that the extension axioms hold almost surely as  $n$  gets large.

# An outrageous conjecture

In their paper *From stability to simplicity* (1998), Kim and Pillay offered the following “outrageous conjecture”:

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This conjecture is false (as we shall see). However, we have:

## Fact (Cherlin-Hrushovski)

*All smoothly approximable theories are simple.*

## Fact (K.)

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$$A \underset{C}{\perp} B \text{ iff } A \cap B \subseteq C$$

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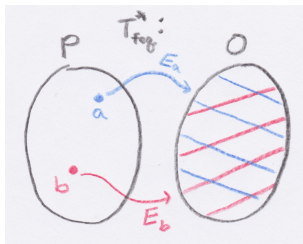
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In fact, a countably categorical theory with disjoint  $n$ -amalgamation for all  $n$  will be  $\omega$ -simple ( $n$ -simple for all  $n$ ) in the sense of Kim, Kolesnikov, and Tsuboi.



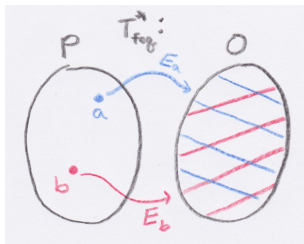
# Counterexample: $T_{\text{feq}}^*$

Our language has two sorts,  $O$  (objects), and  $P$  (parameters), and one ternary relation  $E_x(y, z)$ .  $K_{\text{feq}}$  is the class of finite structures with the property that for all  $a \in P$ ,  $E_a(y, z)$  is an equivalence relation on  $O$ .



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$K_{\text{feq}}$  is a Fraïssé class. Its generic theory,  $T_{\text{feq}}^*$ , is not simple (Shelah) but is NSOP<sub>1</sub> (Chernikov and Ramsey).

### Theorem (K.)

$T_{\text{feq}}^*$  is pseudofinite.

# The uniform measures

Let  $K_{\text{feq}}(n, m)$  be the structures in  $K_{\text{feq}}$  with  $O$ -sort  $[n]$  and  $P$ -sort  $[m]$ , and let  $\mu_{n,m}$  be the uniform measure on  $K_{\text{feq}}(n, m)$ .  $T_{\text{feq}}^*$  is *not* the almost-sure theory of  $(K_{\text{feq}}(n, m), \mu_{n,m})$  as  $n$  and  $m$  grow.

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**Fact (Flajolet and Sedgewick, Proposition VIII.8)**

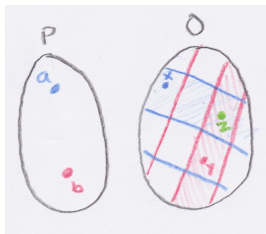
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$T_{\text{feq}}^*$  says:

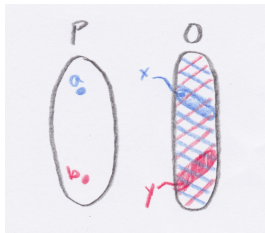
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But asymptotically:

$$\forall(a : P) \forall(b : P) \forall(x : O) \forall(y : O) \exists(z : O) ((a \neq b) \rightarrow E_a(x, z) \wedge E_b(y, z))$$

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## Question

Does  $K_{\text{feq}}$  have a zero-one law for the uniform measures?

## Definition

A Fraïssé class  $K$  is *filtered* by a chain  $K_0 \subseteq K_1 \subseteq K_2 \subseteq \dots$  if each  $K_n$  is a Fraïssé class, and  $\bigcup_{n \in \omega} K_n = K$ .

The generic theory  $T_K$  of  $K$  is the limit of generic theories  $T_{K_n}$ .  
 $\varphi \in T_K$  if and only if  $\varphi \in T_{K_n}$  for all sufficiently large  $n$ .

Hence to show that  $T_K$  is pseudofinite, it suffices to show that each  $T_{K_n}$  is pseudofinite.



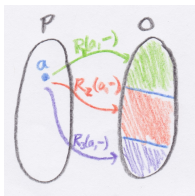
## Filtering $K_{\text{feq}}$

Let  $K_n$  be the subclass of  $K_{\text{feq}}$  consisting of those structures in which the equivalence relation  $E_a$  has at most  $n$  classes, for each parameter  $a \in P$ .

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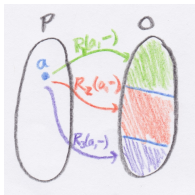
Let  $L_n$  be the expanded language obtained by adding  $n$  binary relations  $R^1(x, y), \dots, R^n(x, y)$ . Then  $K_n$  admits an expansion to a Fraïssé class  $K'_n$ , obtained by naming each  $E_a$ -equivalence class by one of the formulas  $R^i(a, y)$  for each parameter  $a \in P$ .



# Filtering $K_{\text{feq}}$

Let  $K_n$  be the subclass of  $K_{\text{feq}}$  consisting of those structures in which the equivalence relation  $E_a$  has at most  $n$  classes, for each parameter  $a \in P$ .

Let  $L_n$  be the expanded language obtained by adding  $n$  binary relations  $R^1(x, y), \dots, R^n(x, y)$ . Then  $K_n$  admits an expansion to a Fraïssé class  $K'_n$ , obtained by naming each  $E_a$ -equivalence class by one of the formulas  $R^i(a, y)$  for each parameter  $a \in P$ .



Each  $K_n$  has disjoint  $k$ -amalgamation for all  $k$ , and hence its generic theory  $T_{K_n}$  is pseudofinite. Thus  $T_{\text{feq}}^*$  is pseudofinite.

# The Henson graph

We were able to push the “expansion with  $n$ -amalgamation” argument to some unsimple theories via the filtered Fraïssé class method. This method cannot apply to the Henson graph (the Fraïssé limit of the class of all triangle-free graphs).

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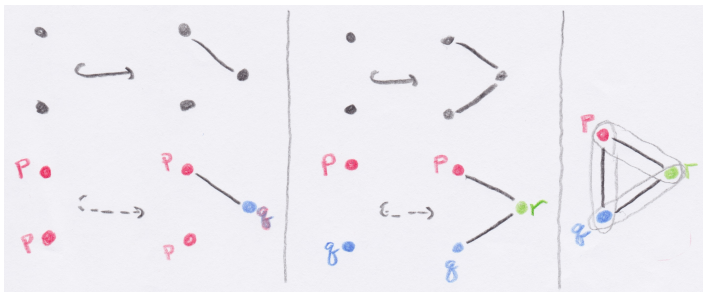


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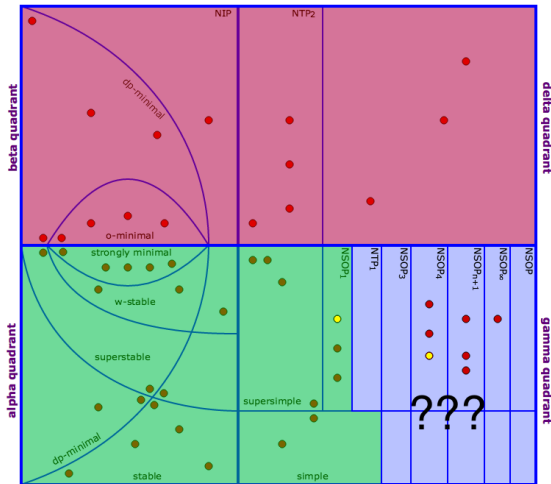
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No Fraïssé class of triangle-free graphs containing the graphs can be expanded to a Fraïssé class with 2- and 3-amalgamation.



# Map of the universe revisited



As far as I know, the question of pseudofiniteness is open for every  $\aleph_0$ -categorical theory in the blue region.

# A new outrageous conjecture

I'd like to boldly modify the conjecture of Kim and Pillay:

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